Unit Commitment: Dealing with the Uncertainties

> Roger J-B Wets University of California, Davis Summer 2014 @ Technical U. Denmark

10 014

Unit commitment Part I



Progressive Hedging: dealing with binary variables

Transmission Network

Figure 1. Topology of the IEEE 300 node system

Transmission Network





NE-ISO net ~30,000 BUS

Figure 1. Topology of the IEEE 300 node system

ISO: Independent System Operator



Federal Energy Regulatory Commission



ISO

In the US is an organization that is responsible for moving electricity over large interstate areas; coordinates, controls and monitors an electricity transmission grid that is larger with much higher voltages than the typical power company's distribution grid.

Is an organization formed at the direction or recommendation of the **FERC**, in the areas where an **ISO** is established, it coordinates, controls and monitors the operation of the electrical power system, usually within a single US State, but sometimes encompassing multiple states.

ISO New England Inc. *(ISO-NE)* is an independent, non-profit RTO, serving Connecticut, Maine, Massachusetts, New Hampshire, Rhode Island and Vermont. Its Board of Directors and its over 400 employees have no financial interest or ties to any company doing business in the region's wholesale electricity marketplace.

Uncertainties

- WEATHER: demand & supply (especially renewables)
- industrial-commercial environment (demand)
- seasonal, day of the week, time of the day
- contingencies: transmission lines, generators



Uncertainties

- WEATHER: demand & supply (especially renewables)
- industrial-commercial environment (demand)
- seasonal, day of the week, time of the day
- contingencies: transmission lines, generators





Uncertainties

- WEATHER: demand & supply (especially renewables)
- industrial-commercial environment (demand)
- seasonal, day of the week, time of the day
- contingencies: transmission lines, generators





Energy Sources



- nuclear energy
- hydro-power
- thermal plants (coal, oil, shale oil, bio, rubish, ...)
- gas turbines (natural gas, from "cracking")
- renewables (wind, solar, ..., ocean waves)

different characteristics







	MISO	NYISO	PJM	ERCOT	CAISO
Market timeline	DA offers due:	DA offers due: 5	DA offers due:	DA bids due	DA offers: 10am
	11am	am	noon	(reserves):	DA results: 1pm
	DA results: 4pm	DA results: 11	DA results: 4pm	1pm/4pm	RT offers: OH -
	Re-bidding due:	am	RT offers due:	DA results	75 min
	5pm	RT offers due:	6pm DA	(reserves):	
	RT offers due:	OH -75 min		1.30pm/6pm	
	OH -30 min			RT offers due:	
				OH -60 min	

Ref: A. Botterud, J. Wang, C. Monteiro, and V. Miranda "Wind Power Forecasting and Electricity Market Operations," available at www.usaee.org/usaee2009/submissions/Onl ineProceedings/Botterud_etal_paper.pdf

between a rock and a hard place



MIP-CPLEX & good ISP-codes (S. Sen & Co.) can only handle effectively problems of moderate size

recall "deadlines"

Our "ARPA-e Team"

- Sandia National Labs: Jean-Paul Watson, César Silva-Monroy, Ross Guttromsom (team builder), John Siirola, William Hart, ...
- * Iowa Sate University: Sarah Ryan, Dinakar Gade, Yonghan Feng, Youngrok Lee
- University of California, Davis: David Woodruff, Roger J-B Wets, Ignacio Rios, Kai Spürkel, Fabian Rüdel, (+ Chuangyin Dang, Julia Peyre ... later this year)
- Alstom: Kwok Cheung (+ ...)
- * @ New-England ISO: Eugene Litvinov (& Joe Mercer, William Callan)
- Unofficial associates: Johannes Royset (NPS), Hoa Chen (UCD) uncertainty design

Minimize $\sum_{k \in K} \sum_{j \in J} c_j^P(k) + c_j^u(k) + c_j^d(k) \quad \text{with}$

$$\sum_{j \in J} p_j(k) = D(k), \quad \forall k \in K$$
$$\sum_{j \in J} \bar{p}_j(k) \ge D(k) + R(k), \quad \forall k \in K$$
$$p_j(k), \bar{p}_j(k) \in \Pi, \quad \forall j \in J, \; \forall k \in K$$

 $\begin{array}{ll} \text{Minimize} \sum_{k \in K} \sum_{j \in J} c_j^P(k) + c_j^u(k) + c_j^d(k) & \text{with} \\ J \text{ generating units} \end{array} \end{array}$

$$\sum_{j \in J} p_j(k) = D(k), \quad \forall k \in K$$
$$\sum_{j \in J} \bar{p}_j(k) \ge D(k) + R(k), \quad \forall k \in K$$
$$p_j(k), \bar{p}_j(k) \in \Pi, \quad \forall j \in J, \; \forall k \in K$$

Minimize $\sum_{k \in K} \sum_{j \in J} c_j^P(k) + c_j^u(k) + c_j^d(k)$ with *K time periods J generating units*

$$\sum_{j \in J} p_j(k) = D(k), \quad \forall k \in K$$
$$\sum_{j \in J} \bar{p}_j(k) \ge D(k) + R(k), \quad \forall k \in K$$
$$p_j(k), \bar{p}_j(k) \in \Pi, \quad \forall j \in J, \; \forall k \in K$$

 $\begin{array}{l} \text{Minimize} & \sum_{k \in K} \sum_{j \in J} c_j^P(k) + c_j^u(k) + c_j^d(k) & \text{with} \\ \text{\textit{K time periods}} & \text{J generating units} \end{array}$

$$\sum_{j \in J} p_j(k) = D(k), \quad \forall k \in K$$
$$\sum_{j \in J} \bar{p}_j(k) \ge D(k) + R(k), \quad \forall k \in K$$
$$p_j(k), \bar{p}_j(k) \in \Pi, \quad \forall j \in J, \; \forall k \in K$$

 $\begin{array}{l} \text{Minimize} & \sum_{k \in K} \sum_{j \in J} c_j^P(k) + c_j^u(k) + c_j^d(k) & \text{with} \\ \textbf{\textit{K time periods}} & \textbf{J generating units} \end{array}$

$$\sum_{j \in J} p_j(k) = D(k), \quad \forall k \in K$$
$$\sum_{j \in J} \bar{p}_j(k) \ge D(k) + R(k), \quad \forall k \in K$$
$$p_j(k), \bar{p}_j(k) \in \Pi, \quad \forall j \in J, \; \forall k \in K$$

Minimize $\sum_{k \in K} \sum_{j \in J} c_j^P(k) + c_j^u(k) + c_j^d(k)$ with K time periods J generating units

$$\sum_{j \in J} p_j(k) = D(k), \quad \forall k \in K$$
$$\sum_{j \in J} \bar{p}_j(k) \ge D(k) + R(k), \quad \forall k \in K$$
$$p_j(k), \bar{p}_j(k) \in \Pi, \quad \forall j \in J, \; \forall k \in K$$

Minimize $\sum_{k \in K} \sum_{j \in J} c_j^P(k) + c_j^u(k) + c_j^d(k)$ with *K* time periods J generating units

$$\begin{split} &\sum_{j \in J} p_j(k) = \frac{\text{demand}}{D(k)}, \ \forall k \in K \\ &\sum_{j \in J} \bar{p}_j(k) \geq D(k) + R(k), \ \forall k \in K \\ &p_j(k), \bar{p}_j(k) \in \Pi, \ \forall j \in J, \ \forall k \in K \end{split}$$

 $\begin{array}{l} \text{production cost startup cost shutdown cost}\\ \text{Minimize} &\sum_{k \in K} \sum_{j \in J} c_j^P(k) + c_j^u(k) + c_j^d(k) \quad \text{with}\\ \text{K time periods} \quad J \text{ generating units} \\ \text{power output} &\sum_{j \in J} p_j(k) = \frac{\text{demand}}{D(k)}, \quad \forall k \in K\\ &\sum_{j \in J} \bar{p}_j(k) \geq D(k) + R(k), \quad \forall k \in K \end{array}$

 $j \in J$ $p_j(k), \overline{p}_j(k) \in \Pi, \quad \forall j \in J, \ \forall k \in K$

 $\begin{array}{l} & \textit{production cost startup cost shutdown cost} \\ & \text{Minimize} \sum_{k \in K} \sum_{j \in J} c_j^P(k) + c_j^u(k) + c_j^d(k) & \text{with} \\ & \textit{K time periods} & J \textit{ generating units} \\ & \textit{power output} \sum_{j \in J} p_j(k) = \frac{\textit{demand}}{D(k)}, \ \forall k \in K \\ & \textit{max power output} \sum_{j \in J} \bar{p}_j(k) \geq D(k) + R(k), \ \forall k \in K \\ & p_j(k), \overline{p}_j(k) \in \Pi, \ \forall j \in J, \ \forall k \in K \end{array}$

 $\begin{array}{l} & \textit{production cost startup cost shutdown cost} \\ & \text{Minimize} \sum_{k \in K} \sum_{j \in J} c_j^P(k) + c_j^u(k) + c_j^d(k) & \text{with} \\ & \textit{K time periods} & J \textit{generating units} \\ & \textit{power output} \sum_{j \in J} p_j(k) = \frac{\textit{demand}}{D(k)}, \ \forall k \in K \\ & \textit{max power output} \sum_{j \in J} \bar{p}_j(k) \geq D(k) + R(k), \ \forall k \in K \\ & p_j(k), \overline{p}_j(k) \in \Pi, \ \forall j \in J, \ \forall k \in K \end{array}$

 $\begin{array}{l} & \textit{production cost startup cost shutdown cost} \\ & \text{Minimize} \sum_{k \in K} \sum_{j \in J} c_j^P(k) + c_j^u(k) + c_j^d(k) & \text{with} \\ & \textit{K time periods} & J \textit{generating units} \\ & \textit{power output} \sum_{j \in J} p_j(k) = \frac{\textit{demand}}{D(k)}, \ \forall k \in K \\ & \textit{max power output} \sum_{j \in J} \bar{p}_j(k) \geq D(k) + R(k), \ \forall k \in K \\ & p_j(k), \overline{p}_j(k) \in \Pi, \ \forall j \in J, \ \forall k \in K \end{array}$

 $\begin{array}{l} \textit{production cost startup cost shutdown cost} \\ \text{Minimize} \sum_{k \in K} \sum_{j \in J} c_j^P(k) + c_j^u(k) + c_j^d(k) \quad \text{with} \\ \textit{K time periods} \quad J \textit{ generating units} \\ \textit{power output} \sum_{j \in J} p_j(k) = \frac{\textit{demand}}{D(k)}, \quad \forall k \in K \\ \textit{max power output} \sum_{j \in J} \bar{p}_j(k) \geq D(k) + R(k), \quad \forall k \in K \\ p_j(k), \bar{p}_j(k) \in \Pi, \quad \forall j \in J, \ \forall k \in K \end{array}$

 Π region of feasible production, all generating units, all time periods. The specific nature of Π is model-dependent.

"Stochastic Version"

min. expectation (actually: risk measure) with penalties

asure) production cost startup cost shutdown cost S Minimize $\sum_{k \in K} \sum_{j \in J} c_j^P(k) + c_j^u(k) + c_j^d(k)$ with K time periods J generating units

power output
$$\sum_{j \in J} p_j(k) = \underbrace{D(k)}_{D(k)}, \ \forall k \in K$$

adjust node balance eq'ns

 $\begin{array}{l} \textit{max power output} \sum_{j \in J} \bar{p}_j(k) \geq D(k) + R(k), \ \forall k \in K \\ p_j(k), \bar{p}_j(k) \in \Pi, \ \forall j \in J, \ \forall k \in K \end{array}$

 Π region of feasible production, all generating units, all time periods. The specific nature of Π is model-dependent.

"Stochastic Version"

Aggregation Principle in Stochastic Optimization

⇒ to Progressive Hedging Algorithm

Here-&-Now vs. Wait-&-See

* Basic Process: decision → observation → decision x¹ → ξ → x²_ξ
* Here-&-Now problem! x¹ not all contingencies can be "protected" by available instruments, i.e.

 Wait-&-See problem: instruments are available to cover all contingencies choose (x¹_ξ, x²_ξ) after observing ξ

Stochastic Optimization: Fundamental Theorem

A here-&-now problem can be transformed in a wait-&-see problem by introducing the

appropriate `contingencies' costs (price of nonanticipativity)

 $\min \mathbb{E}\left\{f_0(\boldsymbol{\xi}, x^1, x_{\boldsymbol{\xi}}^2)\right\}$ $x^1 \in C^1 \subset \mathbb{R}^n$, $x_{\xi}^2 \in C^2(\xi, x^1), \forall \xi.$

Explicit non-anticipativity

 $\min \mathbb{E}\left\{f_0(\boldsymbol{\xi}, x^1, x_{\boldsymbol{\xi}}^2)\right\}$ $x^1 \in C^1 \subset \mathbb{R}^n,$ $x_{\boldsymbol{\xi}}^2 \in C^2(\boldsymbol{\xi}, x^1), \forall \boldsymbol{\xi}.$

 $\min \mathbb{E} \left\{ f_0(\xi, x_{\xi}^1, x_{\xi}^2) \right\}$ $x_{\xi}^1 \in C^1 \subset \mathbb{R}^n,$ $x_{\xi}^2 \in C^2(\xi, x_{\xi}^1), \forall \xi.$

$$\min \mathbb{E}\left\{f_0(\boldsymbol{\xi}, x^1, x_{\boldsymbol{\xi}}^2)\right\}$$
$$x^1 \in C^1 \subset \mathbb{R}^n,$$
$$x_{\boldsymbol{\xi}}^2 \in C^2(\boldsymbol{\xi}, x^1), \forall \boldsymbol{\xi}$$

 $\min \mathbb{E} \left\{ f_0(\xi, x_{\xi}^1, x_{\xi}^2) \right\}$ $x_{\xi}^1 \in C^1 \subset \mathbb{R}^n,$ $x_{\xi}^2 \in C^2(\xi, x_{\xi}^1), \forall \xi.$

$$x_{\xi}^{1} = \mathbb{E}\left\{x_{\xi}^{1}\right\} \quad \forall \xi$$

Explicit non-anticipativity

 $\min \mathbb{E}\left\{f_0(\boldsymbol{\xi}, x^1, x_{\boldsymbol{\xi}}^2)\right\}$ $x^1 \in C^1 \subset \mathbb{R}^n,$ $x_{\boldsymbol{\xi}}^2 \in C^2(\boldsymbol{\xi}, x^1), \forall \boldsymbol{\xi}.$

 $\min \mathbb{E} \left\{ f_0(\xi, x_{\xi}^1, x_{\xi}^2) \right\}$ $x_{\xi}^1 \in C^1 \subset \mathbb{R}^n,$ $x_{\xi}^2 \in C^2(\xi, x_{\xi}^1), \forall \xi.$ $x_{\xi}^1 = \mathbb{E} \left\{ x_{\xi}^1 \right\} \quad \forall \xi$ $w_{\xi} \perp \text{ subspace of constant fcns}$ $multipliers \qquad \Rightarrow \mathbb{E} \left\{ w_{\xi} \right\} = 0$

Explicit non-anticipativity

 $\min \mathbb{E}\left\{f_0(\boldsymbol{\xi}, x^1, x_{\boldsymbol{\xi}}^2)\right\}$ $x^1 \in C^1 \subset \mathbb{R}^n,$ $x_{\boldsymbol{\xi}}^2 \in C^2(\boldsymbol{\xi}, x^1), \forall \boldsymbol{\xi}.$

 $\min \mathbb{E} \left\{ f_0(\xi, x_{\xi}^1, x_{\xi}^2) \right\}$ $x_{\xi}^1 \in C^1 \subset \mathbb{R}^n,$ $x_{\xi}^2 \in C^2(\xi, x_{\xi}^1), \forall \xi.$

 $\begin{aligned} x_{\xi}^{1} &= \mathbb{E}\left\{x_{\xi}^{1}\right\} \quad \forall \xi \\ w_{\xi} \perp \text{ subspace of constant fcns} \\ \textbf{multipliers} \qquad \Rightarrow \mathbb{E}\left\{w_{\xi}\right\} = 0 \\ \text{min } \mathbb{E}\left\{f_{0}(\boldsymbol{\xi}, x_{\xi}^{1}, x_{\xi}^{2}) + \langle w_{\xi}, x_{\xi}^{1} \rangle - \langle w_{\xi}, \mathbb{E}\left\{x_{\xi}^{1}\right\} \rangle\right\} \\ \text{such that } x_{\xi}^{1} \in C_{1}, \quad x_{\xi}^{2} \in C_{2}(\boldsymbol{\xi}, x_{\xi}^{1}) \end{aligned}$

Explicit non-anticipativity

 $\min \mathbb{E}\left\{f_0(\boldsymbol{\xi}, x^1, x_{\boldsymbol{\xi}}^2)\right\}$ $x^1 \in C^1 \subset \mathbb{R}^n,$ $x_{\boldsymbol{\xi}}^2 \in C^2(\boldsymbol{\xi}, x^1), \forall \boldsymbol{\xi}.$

 $\min \mathbb{E} \left\{ f_0(\xi, x_{\xi}^1, x_{\xi}^2) \right\}$ $x_{\xi}^1 \in C^1 \subset \mathbb{R}^n,$ $x_{\xi}^2 \in C^2(\xi, x_{\xi}^1), \forall \xi.$

 $\begin{aligned} x_{\xi}^{1} &= \mathbb{E}\left\{x_{\xi}^{1}\right\} \quad \forall \xi \\ w_{\xi} \perp \text{ subspace of constant fcns} \\ \textbf{multipliers} \qquad \Rightarrow \mathbb{E}\left\{w_{\xi}\right\} = 0 \\ \text{min } \mathbb{E}\left\{f_{0}(\boldsymbol{\xi}, x_{\xi}^{1}, x_{\xi}^{2}) + \langle w_{\xi}, x_{\xi}^{1} \rangle - \langle w_{\xi} \rangle \mathbb{E}\left\{x_{\xi}^{1}\right\} \rangle\right\} \\ \text{such that } x_{\xi}^{1} \in C_{1}, \quad x_{\xi}^{2} \in C_{2}(\boldsymbol{\xi}, x_{\xi}^{1}) \end{aligned}$
Progressive Hedging Algorithm

0. w_{ξ}^{0} such that $\mathbb{E}\left\{w_{\xi}^{0}\right\} = 0$, v = 0. Pick $\rho > 0$ 1. for all ξ : $(x_{\xi}^{1,v}, x_{\xi}^{2,v}) \in \arg\min f_{0}(\xi; x^{1}, x^{2}) - \langle w_{\xi}^{v}, x^{1} \rangle$ $x^{1} \in C^{1} \subset \mathbb{R}^{n_{1}}, x^{2} \in C^{2}(\xi, x^{1}) \subset \mathbb{R}^{n_{2}}$ 2. $\overline{x}^{1,v} = \mathbb{E}\left\{x_{\xi}^{1,v}\right\}$. Stop if $|x_{\xi}^{1,v} - \overline{x}^{1,v}| = 0$ (approx.) otherwise $w_{\xi}^{v+1} = w_{\xi}^{v} + \rho\left[x_{\xi}^{1,v} - \overline{x}^{1,v}\right]$, return to 1. with v = v + 1

Progressive Hedging Algorithm

0.
$$w_{\xi}^{0}$$
 such that $\mathbb{E}\left\{w_{\xi}^{0}\right\} = 0$, $v = 0$. Pick $\rho > 0$
1. for all ξ :
 $(x_{\xi}^{1,v}, x_{\xi}^{2,v}) \in \arg\min f_{0}(\xi; x^{1}, x^{2}) - \langle w_{\xi}^{v}, x^{1} \rangle$
 $x^{1} \in C^{1} \subset \mathbb{R}^{n_{1}}, x^{2} \in C^{2}(\xi, x^{1}) \subset \mathbb{R}^{n_{2}}$
2. $\overline{x}^{1,v} = \mathbb{E}\left\{x_{\xi}^{1,v}\right\}$. Stop if $|x_{\xi}^{1,v} - \overline{x}^{1,v}| = 0$ (approx.)
otherwise $w_{\xi}^{v+1} = w_{\xi}^{v} + \rho\left[x_{\xi}^{1,v} - \overline{x}^{1,v}\right]$, return to 1. with $v = v + 1$

Convergence: add a proximal term

$$f_0(\xi; x^1, x^2) - \langle w_{\xi}^{\nu}, x^1 \rangle - \frac{\rho}{2} |x^1 - \overline{x}^{1,\nu}|^2$$

linear rate in $(x^{1,v}, w^v)$... eminently parallelizable

PH: Implementation issues

implementation: choice of ρ ... scenario (×), *ith*-decision (i) dependent (heuristic) extension to problems with integer variables non-convexities: e.g. ground-water remediation with non-linear PDE recourse

asynchronous

partitioning (= different information feeds) $\min \mathbb{E} \{ f(\xi, x) \}, \quad f(\xi, x) = f_0(x) + \iota_{C(\xi, x)}(x)$ $S = \{ \Xi_1, \Xi_2, \dots, \Xi_K \} \text{ a partitioning of } \Xi, \quad p_k = P(\Xi_k)$ $\mathbb{E} \{ f(\xi, x) \} = \sum_n p_n \mathbb{E} \{ f(\xi, x) | \Xi_n \} \quad \text{(Bundling)}$ $\text{defining } g(k, x) = \mathbb{E} \{ f_0(\xi, x) | \Xi_n \} \text{ if } x \in C_k = \bigcap_{\xi \in \Xi_k} C_{\xi}$ solve the problem as: $\min \sum_{n=1}^N p_k g(k, x)$

Bundling

PH: binary variables

 $\min\langle c, x \rangle + \sum_{\xi \in \Xi} p_{\xi} \langle q_{\xi}, y_{\xi} \rangle$ such that $x \in C_1, \ y_{\xi} \in C_2(\xi, x) \ \forall \xi \in \Xi$ binary (integer) variables: some x's, some y_{ξ} 's.

PH: binary variables

 $\min\langle c, x \rangle + \sum_{\xi \in \Xi} p_{\xi} \langle q_{\xi}, y_{\xi} \rangle \text{ such that} \\ x \in C_1, \ y_{\xi} \in C_2(\xi, x) \ \forall \xi \in \Xi \\ \text{binary (integer) variables: some } x \text{'s, some } y_{\xi} \text{'s.}$

Choice of $\rho \rightarrow \rho_j$ depending on $c_j, |x_j|, ...$ and augmentation

Variable Fixing, in particular binaries, $x_j(s) = \text{constant} (k \text{ iterations})$ Variable Slamming: aggressive variable fixing $x_j(s) \approx \text{constant} (\& c_j x_j(s))$ "Sufficient" variable convergence ~ for small values of $c_j x_j(s)$

Termination criterion: variable slamming when $x_j^{\nu}(\xi) - x_j^{\nu+1}(\xi)$ small

Detecting cycling behavior: (simple) hashing scheme

PH: binary variables

 $\min\langle c, x \rangle + \sum_{\xi \in \Xi} p_{\xi} \langle q_{\xi}, y_{\xi} \rangle \text{ such that} \\ x \in C_1, \ y_{\xi} \in C_2(\xi, x) \ \forall \xi \in \Xi \\ \text{binary (integer) variables: some } x \text{'s, some } y_{\xi} \text{'s.}$

Choice of $\rho \rightarrow \rho_j$ depending on $c_j, |x_j|, ...$ and augmentation

Variable Fixing, in particular binaries, $x_j(s) = \text{constant} (k \text{ iterations})$ Variable Slamming: aggressive variable fixing $x_j(s) \approx \text{constant} (\& c_j x_j(s))$ "Sufficient" variable convergence ~ for small values of $c_j x_j(s)$

Termination criterion: variable slamming when $x_j^{\nu}(\xi) - x_j^{\nu+1}(\xi)$ small

Detecting cycling behavior: (simple) hashing scheme

Enough variables fixed \Rightarrow clean up with CPLEX-MIP

Error Bounds

 $f(\xi, x) = f_0(\xi, x^1, x^2) \text{ if } x^1 \in C^1, x^2 \in C^2(\xi, x^1); +\infty \text{ else}$ Stochastic Program (P):

 $\min_{x \in \mathcal{M}} \mathbb{E} \{ f(\xi, x) \} \text{ such that } x^1 \equiv \mathbb{E} \{ x_{\xi}^1 \}$ Dual Program (*D*)

 $\max_{w \in \mathcal{M}^*} \mathbb{E}\left\{-f^*(\xi, w_{\xi})\right\} \text{ such that } \mathbb{E}\{w_{\xi}\} = 0$ weak duality holds: $\inf P \ge \sup D \Rightarrow \text{ for any feasible } \hat{w}$ $-f^*(\xi, \hat{w}_{\xi}) = \min_x \left[f(\xi, x) + \langle \hat{w}_{\xi}, x^1 \rangle, x \in \mathbb{R}^{n_1 + n_2}\right]$ yields a lower bound for (P), better if \hat{w}_{ξ} is near-optimal \Rightarrow rely on w^* of PH-algorithm to generate lower bound.

Augmentation function

$$m(\Delta, \lambda_{\max}; z, \lambda) = \int_0^{\Delta} \psi(z - s, \lambda, \lambda_{\max}) \varphi(\Delta; s) \, ds, \ z \in [0, \lambda_{\max}]$$



Augmentation function

$$m(\Delta, \lambda_{\max}; z, \lambda) = \int_0^{\Delta} \psi(z - s, \lambda, \lambda_{\max}) \varphi(\Delta; s) \, ds, \ z \in [0, \lambda_{\max}]$$





Unit commitment Part II



Dealing with load (demand) uncertainty

CT load exploratory process

Hourly Load in A Week







CT load vs. weather variables



Robust decisions in a stochastic environment Gemand a robust model of the uncertainty.







... to load on day D

to be delivered-load: l(t)

= $\operatorname{fcn}(\operatorname{temp}(\tau \le t), \operatorname{dewpt}(\tau \le t), \operatorname{clcover}(\tau \le t), \operatorname{wind}(\tau \le t)), t \le 24$



... to load on day D

to be delivered-load: l(t)

= $\operatorname{fcn}(\operatorname{temp}(\tau \le t), \operatorname{dewpt}(\tau \le t), \operatorname{clcover}(\tau \le t), \operatorname{wind}(\tau \le t)), t \le 24$

BUT THAT WOULDN'T CAPTURE THE UNCERTAINTY! ONE WOULD EXPECT:

from predictions on day D-1 to load forecasts on day D



THE DATA













weather prediction @ 11 a.m.
... but too late!

better @ 11 p.m.

weather prediction @ 11 a.m. better @ 11 p.m.
... but too late!

surface wind =>? power wind

weather prediction @ 11 a.m. better @ 11 p.m.
... but too late!

surface wind =>? power wind
 cloud cover (no historical prediction data) -- only
 actuals are available

🛛 weather prediction @ 11 a.m. better @ 11 p.m. ... but too late! surface wind =>? power wind cloud cover (no historical prediction data) -- only actuals are available model to be used for the stochastic load predictions model: SDE, time series, ??? all inappropriate





a) segmentation: season + day characteristics



a) segmentation: season + day characteristicsb) functional regression for given segment



a) segmentation: season + day characteristics
b) functional regression for given segment
c) hourly distribution of errors per segment



a) segmentation: season + day characteristics
 b) functional regression for given segment
 c) hourly distribution of errors per segment
 HOW THIS IS CARRIED OUT ?



a) segmentation: season + day characteristics
 b) functional regression for given segment
 c) hourly distribution of errors per segment
 HOW THIS IS CARRIED OUT ?
 d) conditional distribution of errors => process



a) segmentation: season + day characteristics
 b) functional regression for given segment
 c) hourly distribution of errors per segment
 HOW THIS IS CARRIED OUT ?
 d) conditional distribution of errors => process

e) discretization of the process => scenarios
Segmentation

to enrich ~ similars, analogs (± standard) data: Wednesday rule, zone rule? seasons: (factor analysis, 'heuristics') ± spring & fall : temperature winter: temperature & cloud cover summer: temperature & dew point wind power (at present): handled independently based total load \approx load on stier analogs scenario - wind power scenario











The Regression Problem

find a function *r* that minimizes errors (with respect to $\|\Box\|$) $\sum_{\text{days d in segment}} \sum_{\text{hours h in day}} \left\| r((tmp_{d,h}, hum_{d,h})) - \text{load}_{d,h} \right\|$

an infinite dimensional problem!

Our approach: rely on 2-dimensional epi-splines ("innovation")

- epi-splines approximate with arbitrary accuracy 'any' function
- epi-splines are completely determined by a finite # of parameters
- allows (via constraints) to include 'soft' (non-data) information

The Errors Distributions

Given segment # and associated r, for fixed hour h $e_{d,h} = \text{load}_{d,h} - r((tmp_{d,h}, hum_{d,h})), d \in \text{segment } \#$ \Rightarrow estimate the density f_h of the errors (at h in segment #) yields an overall estimate of the 'volatility' (in fact, more) another infinite dimensional problem & data might be scarce

Our approach: estimation via exponential epi-spline (novel): $-f_h = \exp(-s_h), s_h$ an epi-spline $(\Rightarrow f_h \ge 0)$

- same properties as epi-spline, could include unimodality restriction





Approximation foundation

Function Identification Problem: optim., diff. eqtns, processes, ...

(*FIP*) find $f \in \arg\min\{\psi(f)| f \in F \subset \mathcal{F}\}$

here $\mathcal{F} = \operatorname{lsc-fcns}(\mathbb{R}^n)$ more generally a Polish space \mathcal{F} -approximation: *aw*-topology ~ dl(f,g) epi-distance

Approximation foundation

Function Identification Problem: optim., diff. eqtns, processes, ...

(*FIP*) find $f \in \arg\min\{\psi(f)| f \in F \subset \mathcal{F}\}$

here $\mathcal{F} = \operatorname{lsc-fcns}(\mathbb{R}^n)$ more generally a Polish space \mathcal{F} -approximation: *aw*-topology ~ dl(f,g) epi-distance

 $dl_{\rho}(f,g)$ *g* $dl_{
ho}(f,g)$ d(x, epig) $dl(f,g) = \int dl_{\rho}(f,g)e^{-\rho} d\rho$ d(x, epi)

Epi-splines

 $\mathcal{R} = \{R_1, \dots, R_N \text{ open}\} \text{ partitions (no overlap) } B = \bigcup_{k=1}^N \operatorname{cl} R_k \text{ (closed)} \\ \text{poly}^p(\mathbb{R}^n) \text{ defined by } n_p \leq (n+p)!/(n!p!) \text{ real parameters} \\ \text{lsc epi-spline } s : \mathbb{R}^n \to \overline{\mathbb{R}} \text{ with partition } \mathcal{R} \text{ of order } p \in \mathbb{N}_0 \text{ if} \\ \text{ on each } R_k \ s \in \operatorname{poly}^p(\mathbb{R}^n), \quad s \equiv \infty \text{ on } \mathbb{R}^n \setminus B, \quad s \text{ is lsc on } \mathbb{R}^n \\ \text{ then } s \in \operatorname{e-spl}_n^p(\mathcal{R})..... \subset \operatorname{lsc-fcns}(B) \subset \operatorname{lsc-fcns}(\mathbb{R}^n) \end{cases}$

Epi-splines

 $\mathcal{R} = \{R_1, \dots, R_N \text{ open}\} \text{ partitions (no overlap) } B = \bigcup_{k=1}^N \text{cl} R_k (\text{closed}) \\ \text{poly}^p(\mathbb{R}^n) \text{ defined by } n_p \leq (n+p)!/(n!p!) \text{ real parameters} \\ \text{lsc epi-spline } s : \mathbb{R}^n \to \overline{\mathbb{R}} \text{ with partition } \mathcal{R} \text{ of order } p \in \mathbb{N}_0 \text{ if} \\ \text{ on each } R_k s \in \text{poly}^p(\mathbb{R}^n), \quad s \equiv \infty \text{ on } \mathbb{R}^n \setminus B, \quad s \text{ is lsc on } \mathbb{R}^n \\ \text{ then } s \in \text{e-spl}_n^p(\mathcal{R})..... \subset \text{ lsc-fcns}(B) \subset \text{ lsc-fcns}(\mathbb{R}^n) \\ 1. \forall p \in \mathbb{N}_0, \text{ infnite refinement } \{\mathcal{R}^v\}_{v=1}^{\infty} \text{ of closed } B \subset \mathbb{R}^n \\ \\ \infty \end{bmatrix}$

 $\bigcup_{\nu=1}^{p} e-\operatorname{spl}_{n}^{p}(\mathcal{R}^{\nu}) \text{ is dense in } \operatorname{lsc-fcns}(B)$

refinements: boxes, simplexes, ...

2. When
$$(FIP^{\nu}) \rightarrow_{epi} (FIP), s^k \in \arg\min(FIP^{\nu_k}), dl(s^k, f) \rightarrow 0$$

then $f \in \arg\min(FIP)$.

* curve fitting: *F* known properties of the curve

* curve fitting: *F* known properties of the curve

financial curves: yield curve, discount factor curve

* curve fitting: F known properties of the curve

financial curves: yield curve, discount factor curve

variogram: geostatistics, deposit dispersion

* curve fitting: *F* known properties of the curve

financial curves: yield curve, discount factor curve

variogram: geostatistics, deposit dispersion

uncertainty quantification in a harmonic excitation

- * curve fitting: *F* known properties of the curve
- financial curves: yield curve, discount factor curve
- variogram: geostatistics, deposit dispersion
- uncertainty quantification in a harmonic excitation
- * building stochastic processes: commodities prices, e-loads, ...

- * curve fitting: *F* known properties of the curve
- financial curves: yield curve, discount factor curve
- variogram: geostatistics, deposit dispersion
- uncertainty quantification in a harmonic excitation
- * building stochastic processes: commodities prices, e-loads, ...
- density estimation: ψ max. likelihood (captures observations),
 F soft information: support, shape, bounded moments,
 Bayesian, ..



a. identify all observed load curves in each sub-segment *b*. for each sub-segment: re-calculate regression and errors distribution *c*. repeat for each sub-segment @ (say, 1 p.m.) \Rightarrow sub-sub-segment

Unit commitment Part II



Dealing with renewables-supply uncertainty

Wind & Solar

wind and solar: complementary balance ±

★ wind scenarios: 3TIER analogues (ARPA-e ⊕)

scenario building (state-of-the-art)

 stochastic process building with soft information coming from the dynamics (started)

errors: $e_t^d = a_t^d - f_t^d$ (Bonneville Power Administration) 1. distribution of the forecast for hours of interest 2. segment errors (per day) according to forecast wind power 3. compute conditional error distribution Scenarios: - generate scenario for the errors - generate scenario paths (discrete process)













References

- K. Cheung, Y. Feng, D. Gade, Y. Lee, C. Monroy, I. Rios, F. Rüdel, S. Ryan, J.-P. Watson, R. Wets and D. Woodruff. Stochastic Unit Commitment at ISO scale: an ARPAe projectf. 2014 IEEE Power & Energy Society General Meeting Proceedings, 2014
- K. Cheung, D. Gade, C. Monroy, S. Ryan, J.-P. Watson, R. Wets and D. Woodruff. Toward scalable stochastic unit commitment - Part 2: Assessing solver performance. IEEE Transactions on Power Systems (submitted). 2013
- Y. Feng, I. Rios, S. Ryan, K. Spürkel, J.-P. Watson, R. Wets and D. Woodruff. Toward scalable stochastic unit commitment - Part 1: load scenarios generation. IEEE Transactions on Power Systems (submitted). 2013
- J.-P. Watson and D. Woodruff. Progressive hedging innovations for a class of stochastic mixed integer resource allocation problems, Computational Management Science, 2010+
- R.T. Rockafellar and R. Wets. Nonanticipativity and L1-martingales in stochastic optimization problems. *Mathematical Programming Study*, 6:170–187, 1976
- I. Rios, R. Wets and D. Woodruff, Multi-period forecasting with limited information and scenario generation with limited data. 2013 (submitted for publication)
- J. Royset and R. Wets. Epi-splines and exponential epi-splines: pliable approximation tools, Tech. Report U. of California, Davis, 2013 (submitted)
- J. Royset and R. Wets, Nonparametric density estimation via exponential epi-splines: fusion of soft and hard information. 2013 (submitted)

